

AMBA™ 4 ATB Protocol Specification

ATBv1.0 and ATBv1.1



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Release Information

Change history

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19 June 2006	A	Non-Confidential	First release, AMBA3, ATB version 1.0
28 March 2012	B	Non-Confidential	Second release, AMBA 4, ATB version 1.1

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Glossary

Preface

This preface introduces the *Advanced Microcontroller Bus Architecture* (AMBA 4) *Advanced Trace Bus* (ATB). It contains the following sections:

- *About this book* on page viii
- *Using this book* on page ix
- *Conventions* on page x
- *Additional reading* on page xii
- *Feedback* on page xiii.

About this book

This book describes the AMBA 3 ATB protocol. All references to ATB in this specification refer to AMBA 3 ATB. The information in this document supersedes ATB information located in the *CoreSight Architecture Specification*.

Intended audience

This book is written for the following target audience:

- engineers who are familiar with ARM architecture
- hardware engineers integrating ATB protocol-compliant components into a design
- hardware engineers designing ATB protocol-compliant components
- advanced users of development tools that provide support for ATB protocol functionality.

This book does not document the behavior of individual components unless they form a fundamental part of the architecture.

Using this book

This book is organized into the following chapters:

Chapter 1 *Introduction*

Read this chapter for an overview of the protocol and its interface.

Chapter 2 *Signal Descriptions*

Read this chapter for a description of the protocol signal descriptions. It contains information about naming conventions, signal descriptions and values.

Chapter 3 *Flow Control*

Read this chapter for a description of ATB protocol flow control. This chapter includes information about the connection of an ATB master to an ATB slave, and the use of flow control signals to control data transfer between the master and slave.

Chapter 4 *Other ATB Features*

Read this chapter for a description of the other features of an ATB implementation. It describes the use of *Trace Data IDs* (ATIDs), the flush process for the protocol, and trace trigger and synchronization signaling added in ATBv1.1.

Chapter 5 *Interface Signaling Rules*

Read this chapter for a description of the protocol signal rules. It contains a listing of rules that govern the protocol.

Appendix A *Revisions*

Read this for a description of the technical changes between released issues of this book.

Glossary

Read this for definitions of some terms used in this book. The glossary does not contain terms that are industry standard unless the ARM meaning differs from the generally accepted meaning.

———— **Note** ————

ARM publishes a single glossary that relates to most ARM products, see the *ARM Glossary*, <http://infocenter.arm.com/help/topic/com.arm.doc.aeg0014-/index.html>. A definition in the glossary in this book might be more detailed than the corresponding definition in the *ARM Glossary*.

Conventions

The following sections describe conventions that this book can use:

- [Typographic conventions](#)
- [Signal naming](#)
- [Timing diagrams on page xi](#)
- [Numbers on page xi](#)

Typographic conventions

The typographical conventions are:

<i>italic</i>	Introduces special terminology, denotes internal cross-references and citations, or highlights an important note.
bold	Denotes signal names, and is used for terms in descriptive lists, where appropriate.
monospace	Used for assembler syntax descriptions, pseudocode, and source code examples. Also used in the main text for instruction mnemonics and for references to other items appearing in assembler syntax descriptions, pseudocode, and source code examples.

SMALL CAPITALS

Used for a few terms that have specific technical meanings, and are included in the [Glossary on page Glossary-43](#).

Colored text	Indicates a link. This can be: <ul style="list-style-type: none">• a URL, for example, http://infocenter.arm.com• a cross-reference, that includes the page number of the referenced information if it is not on the current page, for example, Numbers on page xi• a link, to a chapter or appendix, or to a glossary entry, or to the section of the document that defines the colored term, for example Timing diagrams on page xi or Numbers on page xi.
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Signal naming

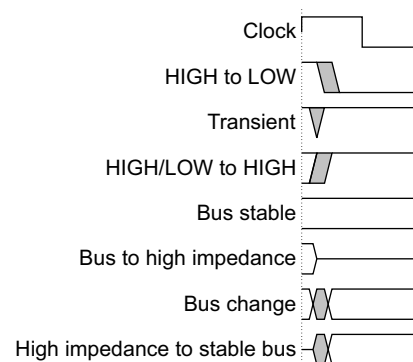
The level of an asserted signal depends on whether the signal is active-HIGH or active-LOW. Asserted means HIGH for active-HIGH signals and LOW for active-LOW signals:

Prefix A	Denotes AMBA 3 <i>Advanced eXtensible Interface</i> (AXI) interface global and address channel signals.
Prefix AF	Denotes signals relating to ATB flush control.
Prefix AT	Denotes signals relating to ATB data flow.
Prefix B	Denotes AXI interface write response channel signals.
Prefix C	Denotes AXI interface low-power interface signals.
Prefix H	Denotes AMBA <i>Advanced High-performance Bus</i> (AHB) signals.
Prefix n	Denotes active-LOW signals except in the case of AHB, <i>Advanced Peripheral Bus</i> (APB), or ATB interface reset signals. These are named HRESETn , PRESETn , and ATRESETn respectively.
Prefix P	Denotes an APB interface signal.
Prefix R	Denotes AXI interface read channel signals.
Prefix W	Denotes AXI interface write channel signals.

Timing diagrams

The figure named *Key to timing diagram conventions* explains the components used in timing diagrams. Variations, when they occur, have clear labels. You must not assume any timing information that is not explicit in the diagrams.

Shaded bus and signal areas are undefined, so the bus or signal can assume any value within the shaded area at that time. The actual level is unimportant and does not affect normal operation.



Key to timing diagram conventions

Timing diagrams sometimes show single-bit signals as HIGH and LOW at the same time and they look similar to the bus change shown in *Key to timing diagram conventions*. If a timing diagram shows a single-bit signal in this way then its value does not affect the accompanying description.

Numbers

Numbers are normally written in decimal. Binary numbers are preceded by `0b`, and hexadecimal numbers by `0x`. In both cases, the prefix and the associated value are written in a monospace font, for example `0xFFFF0000`.

Additional reading

This section lists relevant publications from ARM.

ARM publications

This book contains information that is specific to the ATB protocol. See the following documents for other relevant information:

- *AMBA AXI Protocol Specification AXI3, AXI4, and AXI4-Lite, ACE and ACE-Lite* (ARM IHI 0022).
- *AMBA 3 AHB-Lite Protocol Specification* (ARM IHI 0033)
- *AMBA APB Protocol Specification* (ARM IHI 0024)

Feedback

ARM welcomes feedback on its documentation.

Feedback on this book

If you have comments on the content of this book, send an e-mail to errata@arm.com. Give:

- the title
- the number, ARM IHI 0032B
- the page numbers to which your comments apply
- a concise explanation of your comments.

ARM also welcomes general suggestions for additions and improvements.

———— **Note** ————

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Chapter 1

Introduction

This chapter introduces the ATB protocol. It contains the following sections:

- [About the ATB protocol on page 1-16](#)
- [About the ATB interface on page 1-17.](#)

1.1 About the ATB protocol

The ATB protocol is part of the AMBA protocol family:

- ATBv1.0 is defined as part of AMBA 3
- ATBv1.1 is defined as part of AMBA 4.

The ATB protocol defines how trace information transfers between components in a trace system. ATB is a common bus used by trace components to pass format-independent trace data through a CoreSight system.

A trace component or platform that has trace capabilities requires an ATB interface. The ATB interfaces are designated by function as one of:

Master An interface that generates trace data onto the ATB bus.

Slave An interface that receives trace data from the ATB bus.

The ATB interface supports various features, including:

- stalling of data, using valid and ready responses
- control signals that indicate the number of bytes valid in a cycle
- identification of the originating component, by signaling an associated ID with each data packet
- support for any trace protocol information, data information or data format requirements
- identification of data from all originating components
- flushing.

1.2 About the ATB interface

The ATB is a common bus used by the trace components to pass trace data in a system in a data-agnostic format. The ATB protocol defines the bus behavior and the interface signals are named according to their function.

Figure 1-1 shows the general relationships between the ATB and associated components.

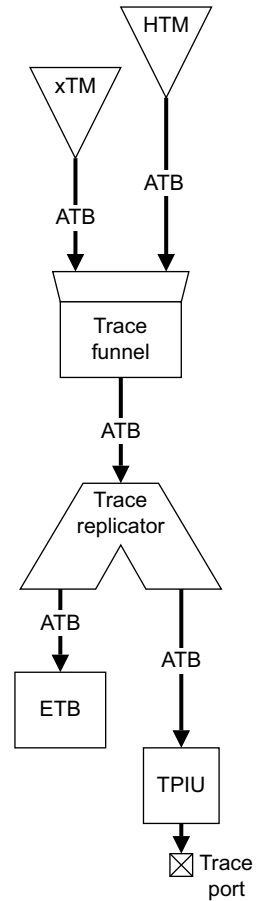


Figure 1-1 ATB relationships

The major components shown in Figure 1-1 are:

Embedded Trace Buffer (ETB)

This is an on-chip trace data storage device.

Trace Macrocell (xTM)

This generates and outputs trace information from a connected processor. ARM trace macrocells include:

- The *Embedded Trace Macrocell* (ETM). An ETM generates instruction trace, and might generate data trace.
- The *Program Flow Trace Macrocell* (PTM). A PTM generates instruction trace.

Advanced High-performance bus Trace Macrocell (HTM)

This outputs trace information describing the AHB interconnect it is monitoring.

Trace Port Interface Unit (TPIU)

This is a device that connects ATB to a trace port, to provide external storage of trace data.

Trace Replicator

This replicates a single ATB slave interface into two independent ATB masters

Trace funnel This combines multiple trace sources onto a single bus.

A trace system includes:

- One or more sources of trace data, for example ETM or HTM.
- One or more trace sinks, or receivers of the data. In [Figure 1-1 on page 1-17](#) these sinks are ETB and TPIU.

Between the trace sources and the sinks, there can be trace links that help manage the passage of data over ATB interfaces, such as the trace funnel and trace replicator.

The trace sources generate ATB data traffic that can be managed by trace links on the ATB path. Trace links can re-transmit information to the trace sinks. Trace sinks act as the receivers of the trace data generated by the trace sources.

Typically a trace sink can request a flush of the trace structure to ensure all trace data, up to a particular time, has been received. These flush requests propagate back up the ATB system from the trace sinks towards the trace sources. On receiving a flush request, a trace source marks all data that occurred before the flush request, and indicates completion when it has output all of this data. Any trace link in the ATB system must:

- pass flush requests up to its connected trace sources
- ensure that it acknowledges flush completion only after all connected trace sources have acknowledged completion.

Chapter 2

Signal Descriptions

This chapter describes the ATB interface signals and signaling rules. It contains the following section:

- [ATB signals on page 2-20.](#)

2.1 ATB signals

This following sections describe the ATB interface signals, by functional group, and any signaling rules:

- [Global signals](#)
- [Data signals](#)
- [Flush control signals on page 2-21](#)
- [Synchronization request signals, ATBv1.1 only on page 2-21.](#)

2.1.1 Global signals

[Table 2-1](#) shows the global signals for a component implementing an ATB trace data interface.

Table 2-1 Global signals

Name	Master	Slave	Description
ATCLK	Input	Input	Global ATB clock.
ATCLKEN	Input	Input	Enable signal for ATCLK domain. From ATBv1.1, this signal is optional.
ATRESETn	Input	Input	ATB interface reset, active LOW. This signal is asserted LOW asynchronously, and deasserted HIGH synchronously.

Clock and reset signals apply to both AT and AF types.

2.1.2 Data signals

[Table 2-2](#) shows the data signals for a component implementing an ATB trace data interface. The table shows the required clamp value of each signal. This is the value required when the component is powered down or disabled.

Table 2-2 Data signals

Name	Master	Slave	Clamp value	Description
ATBYTES[m:0] ^a	Output	Input	LOW	The number of bytes on ATDATA to be captured, minus 1.
ATDATA[n:0] ^a	Output	Input	LOW	Trace data.
ATID[6:0]	Output	Input	LOW	An ID that uniquely identifies the source of the trace. See ATIDs on page 4-30 .
ATREADY	Input	Output	HIGH	Slave is ready to accept data. See Chapter 3 Flow Control .
ATVALID	Output	Input	LOW	A transfer is valid during this cycle. If LOW, all other AT signals must be ignored. See Chapter 3 Flow Control .

a. [Relationship between ATDATA and ATBYTES on page 3-26](#) explains the relationship between the values of m and n.

2.1.3 Flush control signals

Table 2-3 shows the flush control signals for a component implementing an ATB trace data interface. The table shows the required clamp value of each signal. This is the value required when the component is powered down or disabled.

Table 2-3 Flush control signals

Name	Master	Slave	Clamp value	Description
AFVALID	Input	Output	LOW	The flush signal. It is asserted to indicate that all buffers must be flushed because trace capture is about to stop. See Buffer flush on page 4-31 .
AFREADY	Output	Input	HIGH	The flush acknowledge. It is asserted to indicate that buffers have been flushed. See Buffer flush on page 4-31 .

2.1.4 Synchronization request signals, ATBv1.1 only

Table 2-4 shows the recommended synchronization request signal on an ATBv1.1 implementation. The table shows the required clamp value of the signal. This is the value required when the component is powered down or disabled.

Table 2-4 Synchronization request signals

Name	Master	Slave	Clamp value	Description
SYNCREQ	Input	Output	LOW	The synchronization request signal. It is asserted to request the insertion of synchronization information in the trace stream. See Buffer flush on page 4-31 .

Chapter 3

Flow Control

This chapter shows the interconnection of an ATB master and an ATB slave, and describes the ATB protocol flow control, that uses the **ATVALID** and **ATREADY** signals. It contains the following sections:

- [About flow control on page 3-24](#)
- [ATB connections and flow control signaling on page 3-25.](#)

3.1 About flow control

The ATB flow control mechanism is a two-way flow control that the master and the slave can use to control the rate at which the data and control information moves. The source generates the **ATVALID** signal to indicate when the data or control information is available. The destination generates the **ATREADY** signal to indicate that it accepts the data or control information. Transfer of data occurs only when both the **ATVALID** and **ATREADY** signals are HIGH. A specific ID can be assigned to each data item in a cycle, to identify data streams.

3.2 ATB connections and flow control signaling

Figure 3-1 shows the interconnection of an ATB master and an ATB slave.

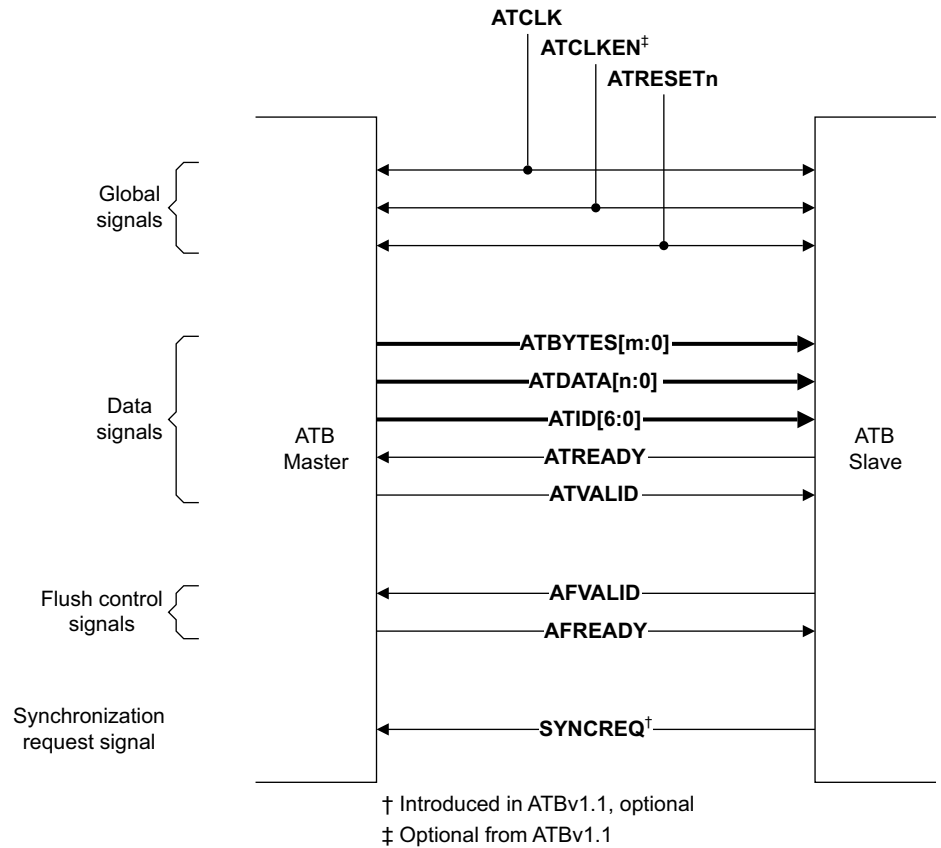


Figure 3-1 ATB signaling between master and slave

ATVALID and **ATREADY** provide a handshake between the master and the slave, to indicate trace data item availability:

- The master indicates that it has trace available to output by asserting **ATVALID**.
- If the slave can accept trace data, it responds by asserting **ATREADY**.
If the slave cannot accept trace data, or does not assert the **ATREADY** signal, the same trace must be output again on the next cycle. [Figure 3-2 on page 3-26](#) shows this process.

A trace source can only start driving **ATVALID** HIGH after **ATRESETn** has been HIGH at a rising edge of **ATCLK**.

Although **ATRESETn** can be asserted LOW asynchronously, deassertion must be synchronous with a rising edge of **ATCLK**.

If **ATREADY** is LOW and **ATVALID** is HIGH on a cycle, **ATVALID**, **ATDATA**, **ATID**, and **ATBYTES** must retain the same value on the next cycle.

If **ATVALID** is LOW, **ATREADY** must be ignored.

[Figure 3-2 on page 3-26](#) shows the **ATVALID** and **ATREADY** flow control.

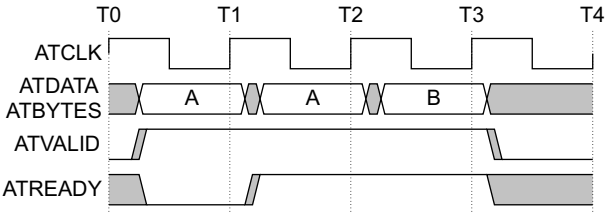


Figure 3-2 ATVALID and ATREADY flow control signaling

Table 3-1 shows the status indicated by the ATB signals at each of the ATCLK rising edges, T1-T4, in Figure 3-2.

Table 3-1 ATB status

Clock cycle	State
T1	Stalled, ATREADY not asserted
T2	A accepted
T3	B accepted
T4	Ignored, not valid

If an implementation of an ATB slave cannot guarantee that it can respond with ATREADY during the same cycle that ATVALID is asserted, then this protocol recommends that the slave implements sufficient internal buffering. This internal buffering must be sufficient to store one or more cycles of trace. The slave can then assert ATREADY when space is available in the buffer, even when ATVALID is not asserted.

The following sections describe the relationship between ATDATA and ATBYTES, and what happens if a master or slave cannot respond:

- Relationship between ATDATA and ATBYTES
- Slave unable to respond on page 3-27
- Master unable to respond on page 3-27.

3.2.1 Relationship between ATDATA and ATBYTES

The following equation defines the relationship between the widths of ATBYTES and ATDATA, defined by the values of m and n in ATBYTES[m:0] and ATDATA[n:0]:

m = log2(n+1) - 4

Using this equation, Table 3-2 shows the relationship between ATBYTES[m:0] and ATDATA[n:0].

Table 3-2 Width relationship between ATDATA and ATBYTES

ATDATA[n:0]	ATBYTES[m:0]
n = 7	ATBYTES not required
n = 15	m = 0
n = 31	m = 1
n = 63	m = 2
n = 127	m = 3

For example, if ATDATA is 32 bits wide (n=31), then ATBYTES is 2 bits wide (m=1).

If **ATREADY** and **ATVALID** are both HIGH, then the bottom bytes of **ATDATA** must be captured and the data must be aligned to the least significant bit.

Note

- Zero bytes cannot be captured.
 - The width of **ATDATA** must be a power of two equal to or greater than eight bits.
-

3.2.2 Slave unable to respond

If an ATB slave interface cannot respond then it must drive **ATREADY** HIGH and **AFVALID** LOW.

Note

This ensures a replicator does not block because one of its two outputs is disabled.

Examples of when a slave cannot respond include:

- The slave is powered down.
- A slave is not present.
- The system has been programmed wrongly.

3.2.3 Master unable to respond

If an ATB master interface cannot respond, then it must drive **AFREADY** HIGH and **ATVALID** LOW.

Examples of when a master cannot respond include:

- The master is powered down.
- A master is not present
- The system has been programmed wrongly.

Chapter 4

Other ATB Features

This chapter describes the features of the ATB specification that are outside the scope of [Chapter 2 Signal Descriptions](#) and [Chapter 3 Flow Control](#). It contains the following sections:

- [ATIDs on page 4-30](#)
- [Buffer flush on page 4-31](#)
- [Signaling a trace trigger, ATBv1.1 on page 4-34](#)
- [Synchronization request, ATBv1.1 on page 4-35.](#)

4.1 ATIDs

Trace data items are generated with separate IDs. These separate IDs can provide:

- differentiation between trace from different sources
- identification of high bandwidth and low bandwidth components of a trace, so components downstream from the trace source can perform selective filtering
- alignment of synchronization information by changing the ID at an alignment synchronization point, such as the beginning of a trace packet.

Most sources use a single, static ID.

ATB trace data items can use:

- any ATID values in the range 0x01-0x6F
- in ATBv1.1, ATID value 0x7D, but only to generate a trace trigger, see [Signaling a trace trigger, ATBv1.1 on page 4-34](#).

Note

The **ATID** values 0x00, 0x70-0x7C, 0x7E, and 0x7F are reserved for use by the CoreSight Architecture and must not be used as ATB IDs.

Because the ID for each trace stream must be unique, there are two options:

Fixed IDs The IDs for all ATB interface sources are chosen when designing the system, and no ATB interfaces are exported from the system.

Programmed IDs

Because the ID for each ATB interface source is programmable by the debugger, components can be reused in larger systems.

Reusable CoreSight components must implement Programmed IDs.

Implementations must ensure that the **ATDATA** value is captured at the same time as bytes on **ATID** are captured.

4.1.1 Trace byte order

An ATB implementation must preserve the order of trace bytes, even between trace with different **ATIDs**.

Although a CoreSight trace funnel can order trace from different sources in any order, when funneled onto a single bus the order cannot be changed, see [Buffer flush on page 4-31](#).

Note

This differs from the required ordering by ID signals in AXI interfaces, where ordering is preserved only between transactions with the same ID.

4.2 Buffer flush

The characteristics of the buffering of trace data items mean it is often necessary to remove remaining data from the buffer to prepare for new data as cycles progress. This process of removing data from a buffer is called flushing.

In a typical trace source there is a fixed amount of time and number of pipeline stages between the occurrence of an event to be traced, and the trace generation for that event. An example of this is an instruction executed by the processor, [Figure 4-1](#) shows an example of a trace generation.

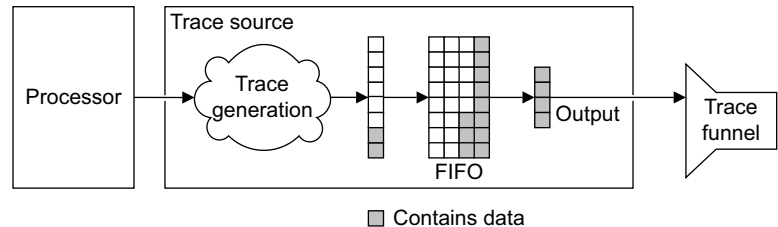


Figure 4-1 Trace generation and output

When a trace data item is generated, it is written to a FIFO. Data is output from the FIFO a whole dataword at a time, where the width of the **ATDATA** bus defines the size of the dataword.

The period of time between the trace generation and output is, in theory, unlimited, because the FIFO might never receive enough trace data to fill a whole output dataword. An example of a possibly extended delay between generation and output is when the trace source withholds output until it has a whole dataword of trace data available for output.

In a system where data is output only when a whole dataword is available, there are various situations that might require a buffer flush, for example:

- The system, or a portion of the system, is about to be powered down or its clock is about to be stopped.
Any trace remaining within buffers or FIFOs in trace sources must be output before powering down or stopping the clock. Therefore, the **AFVALID** signal signals a flush. Normally, after signaling the flush, no more trace is generated because processor and memory activity has stopped. If trace is generated after the flush it can be ignored. An example of trace data generated after a flush is that from a processor idle loop.
- The trace capture device is about to stop capturing, usually because of a trigger point.
The trace capture device might be an off-chip *Trace Port Analyzer* (TPA) or an on-chip *Embedded Trace Buffer* (ETB).

In this case the possibilities include:

- The on-chip logic is signaled that capture is about to stop.
A flush is issued at this point.
- The on-chip logic is signaled about the trigger but does not have any information about how much additional trace is to be captured.
A flush can be issued at the point of the trigger. This ensures all trace generated before the trigger is captured, but makes no difference to trace generated after the trigger.
- A flush is issued periodically.
The period defined for the flush must be chosen so that a flush always occurs between a trigger occurring and the trace capture stopping.
When a flush occurs, indicated by **AFVALID** HIGH, the protocol expects that trace funnels give the highest priority to trace sources that have not yet asserted **AFREADY**.

Note

- A flush sequence can not be cancelled by a slave interface after it begins. This means it cannot be cancelled by the trace sink.
- An implementation must ensure that **ATB** signals are sampled on the rising edge of **ATCLK**, and that **AFVALID** remains asserted until **AFREADY** is deasserted.

AFREADY is asserted when all trace generated at the time **AFVALID** was asserted has been output. This assertion of **AFREADY** does not indicate that the FIFO is empty.

Trace that is generated after **AFVALID** is asserted is stored in the FIFO and is output after **AFREADY** is asserted.

When **AFVALID** is asserted, the ATB master must start outputting buffered trace immediately.

When the ATB master asserts **AFREADY**, then the ATB slave must deassert **AFVALID** on the following cycle, unless an additional flush is required.

When **AFVALID** is asserted, then the master must assert **AFREADY** one cycle after it has output the data that was buffered on the cycle in which **AFVALID** was first asserted. That is, it must assert **AFREADY** one cycle after it has output the last of the data that the assertion of **AFVALID** required it to flush. For more information, see [Buffer flush on page 4-31](#).

Figure 4-2 shows an example of an ATB flush.

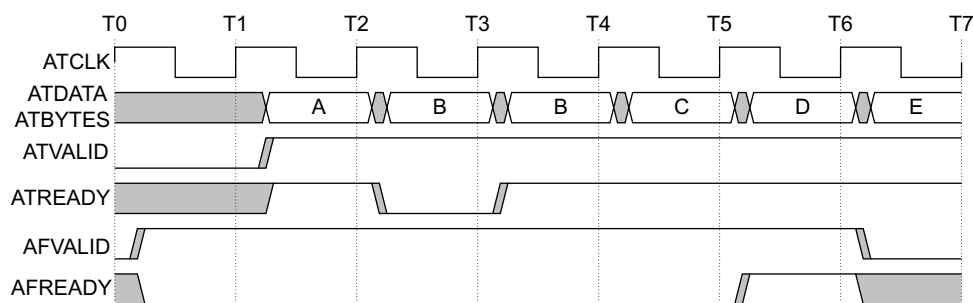


Figure 4-2 Flush signaling

Table 4-1 describes the events that occur during the ATB flush shown in Figure 4-2.

Table 4-1 Flush events

Clock cycle	Event
T1-T2	Flush requested. A, B, and C are held in the FIFO.
T2	Draining begins.
T3-T5	Stalls are still possible.
T6	Flush complete. Output continues.
T6	AFVALID deasserted.

4.2.1 Behavior of trace sources that do not store trace locally

If an ATB master interface is a trace source that does not store any trace locally, [Table 4-2](#) shows how it controls **AFREADY**.

Table 4-2 Basic **AFREADY** control algorithm

ATVALID	ATREADY	AFREADY next cycle
0	-	1
1	0	0
1	1	1

Figure 4-3 on page 4-33 shows this **AFREADY** control for a master that does not store any trace locally.

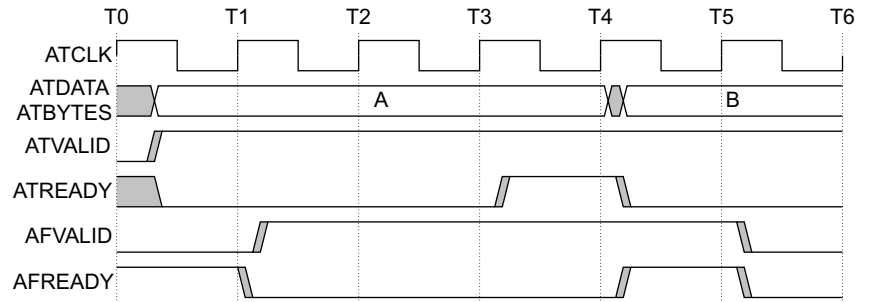


Figure 4-3 Flushing from a master that lacks internal storage

Table 4-3 describes the events, shown in Figure 4-3, that occur during the flush from a master that does not store any trace locally.

Table 4-3 Events associated with flushing from a master that does not store trace locally

Clock cycle	Event
T1	Data (A) generated in the master
T2	Flush request, data (A) is now old data
T3	Flush is not completed because old data (A) is still present in the master
T4	Data (A) is accepted by the slave
T5-T6	New data (B) is generated in the master, flush request acknowledged

4.3 Signaling a trace trigger, ATBv1.1

In ATBv1.1, an ATB master can signal a trace trigger as follows:

- The **ATID** value must be 0x7D.
- **ATBYTES** indicates the number of trace triggers output in this cycle. This protocol specification recommends that, wherever possible, only one trace trigger is generated at a time. In this case:
 - there is only one byte of data
 - **ATBYTES** is zero.
- **ATDATA** indicates the source of each signaled trace trigger, as follows:
 - Each byte of data indicates a separate trace trigger.
 - A byte of data that is zero indicates a trigger from an unknown source. This protocol specification recommends that trace sources do not use this value.
 - A byte of data that is non-zero indicates the ID of the trace source that generated the trace trigger.

For example:

- A trace source that uses ATID 0x10 for its ATB signaling can indicate a trigger on its ATB interface by generating a byte of trace data with:
 - **ATID** signaling 0x7D, indicating a trace trigger
 - **ATBYTES** signaling 0x0, indicating a single data byte
 - **ATDATA** signaling 0x10, indicating the trace source ID.
- A trace stream that uses ATIDs 0x10 and 0x12 can indicate simultaneous trace triggers on both trace sources by generating two bytes of trace data with:
 - **ATID** signaling 0x7D, indicating a trace trigger
 - **ATBYTES** signaling 0x1, indicating two data bytes
 - **ATDATA** signaling 0x1210, indicating the trace source IDs of 0x10 and 0x12.

A trace trigger signal is an IMPLEMENTATION DEFINED message from a trace source to the trace sink. Typically, it might be used to indicate that a significant condition has occurred, and that trace capture should prepare to stop.

4.4 Synchronization request, ATBv1.1

Optionally, an ATBv1.1 implementation can include the **SYNCREQ** synchronization request signal, see [Synchronization request signals, ATBv1.1 only on page 2-21](#).

An ATB slave signals a synchronization request to an ATB master by asserting **SYNCREQ** HIGH for a single **ATCLK** cycle.

Optionally, an ATB master can implement a **SYNCREQ** input, and recognize a synchronization request from an ATB slave when this input is asserted HIGH for a single **ATCLK** cycle.

The **SYNCREQ** signal is not related in any way to any other signal on the ATB interface. However, in an ATBv1.1 implementation, the specification expects:

- **SYNCREQ** to be routed through a system with the other ATB signals
- any trace link, such as a trace funnel, to pass an input synchronization request from an ATB master interface to an upstream trace sources.

An enabled trace source must attempt to provide synchronisation information in the trace stream once for each synchronisation request received. The trace source can:

- delay outputting the synchronisation information to avoid overflowing any internal buffers
- ignore a synchronisation request that is received before the trace source has output the synchronisation information corresponding to a previous synchronisation request.

Chapter 5

Interface Signaling Rules

This chapter describes the ATB interface signaling rules. It contains the following section:

- [Interface signaling rules on page 5-38.](#)

5.1 Interface signaling rules

ATB protocol signaling must conform to the following rules:

- ATB signals must be sampled on the rising edge of **ATCLK**.
- If **ATREADY** is LOW and **ATVALID** is HIGH on a cycle, **ATVALID**, **ATDATA**, **ATID**, and **ATBYTES** must retain the same value on the next cycle.
If **ATVALID** is LOW, **ATREADY** must be ignored.
- If **ATREADY** and **ATVALID** are HIGH, the bottom bytes of **ATDATA** must be captured.
Data must be aligned to the least significant byte.

———— **Note** ————

Zero bytes cannot be captured.

- The width of **ATDATA**, in bits, must be a power of two equal to or greater than eight.
- The relationship between the widths of **ATBYTES[m:0]** and **ATDATA[n:0]** is defined by the equation:
$$m = \log_2(n+1) - 4$$

For examples of possible implementations see [Relationship between ATDATA and ATBYTES on page 3-26](#).
- The **ATDATA** value must be captured at the same time as bytes on **ATID** are captured.
- The order of trace bytes must be preserved, even between trace with different **ATIDs**.
A CoreSight trace funnel can order trace from different sources in any order, but when funneled onto a single bus, the order cannot be changed, see [Buffer flush on page 4-31](#).

———— **Note** ————

This differs from the ordering required by AXI interfaces, where ordering is preserved only between transactions with the same ID.

- When **AFVALID** is asserted, it must remain asserted until **AFREADY** is asserted.
AFREADY is ignored while **AFVALID** is deasserted.
- When **AFVALID** is asserted, an ATB master must output any buffered trace immediately.
- When **AFREADY** is asserted, **AFVALID** must be deasserted on the following cycle, unless an additional flush is required.
- When **AFVALID** is asserted, the ATB master must assert **AFREADY** one cycle after it has output the last of the trace that it generated and stored, up to and including any trace generated on the cycle in which **AFVALID** was asserted. For more information, see [Buffer flush on page 4-31](#).
- **ATRESETn** can be asserted asynchronously. However, **ATRESETn** must be deasserted synchronously on a rising edge of **ATCLK**.

———— **Note** ————

ATRESETn is an active LOW signal, meaning it must be asserted LOW.

- A trace source can only start driving **ATVALID** HIGH after **ATRESETn** has been HIGH at a rising edge of **ATCLK**.
- Whenever an ATB slave interface cannot respond, **ATREADY** must be driven HIGH and **AFVALID** must be driven LOW. Examples of when this must happen include powering down, the slave not being present, and the result of incorrect programming of the system.
This ensures a replicator does not block because one of its two outputs are disabled.

- Whenever an ATB master interface cannot respond, **AFREADY** must be driven HIGH and **ATVALID** must be driven LOW. Examples of when this must happen include powering down, the master not being present, and the result of incorrect programming of the system.
- **ATID**, **ATBYTES**, and **ATDATA** can have any value. This permitted behavior can be used to tie off unused CoreSight trace funnel slave signals.
- The **ATID** values 0x00, 0x70-0x7C, 0x7E, and 0x7F are reserved for use by the CoreSight Architecture and must not be used as ATB IDs.

Appendix A

Revisions

This appendix describes the technical changes between released issues of this book.

Table A-1 Differences between issue A and issue B

Change	Location	Affects
Addition of SYNCREQ signal	<ul style="list-style-type: none">Synchronization request signals, ATBv1.1 only on page 2-21Figure 3-1 on page 3-25	v1.1
Change to reserved ATID values	<ul style="list-style-type: none">ATIDs on page 4-30Interface signaling rules on page 5-38	v1.1
Addition of trace synchronization signaling	Signaling a trace trigger, ATBv1.1 on page 4-34	v1.1
Addition of synchronization request signaling	Synchronization request, ATBv1.1 on page 4-35	v1.1
ATCLKEN becomes an optional signal	<ul style="list-style-type: none">Global signals on page 2-20Figure 3-1 on page 3-25	v1.1
Removal of recommendation that ATDATA is at least 32 bits wide	<ul style="list-style-type: none">Relationship between ATDATA and ATBYTES on page 3-26Interface signaling rules on page 5-38	All versions

Glossary

This glossary describes some of the terms used in technical documents from ARM Limited.

Advanced eXtensible Interface (AXI)

A bus protocol that supports separate address/control and data phases, unaligned data transfers using byte strobes, burst-based transactions with only start address issued, separate read and write data channels to enable low-cost DMA, ability to issue multiple outstanding addresses, out-of-order transaction completion, and easy addition of register stages to provide timing closure. The AXI protocol also includes optional extensions to cover signaling for low-power operation.

AXI is targeted at high performance, high clock frequency system designs and includes a number of features that make it very suitable for high speed sub-micron interconnect.

Advanced High-performance Bus (AHB)

A bus protocol with a fixed pipeline between address/control and data phases. It only supports a subset of the functionality provided by the AMBA AXI protocol. The full AMBA AHB protocol specification includes a number of features that are not commonly required for master and slave IP developments and ARM Limited recommends only a subset of the protocol is usually used. This subset is defined as the AMBA AHB-Lite protocol.

See also Advanced Microcontroller Bus Architecture.

Advanced Microcontroller Bus Architecture (AMBA)

A family of protocol specifications that describe a strategy for the interconnect. AMBA is the ARM open standard for on-chip buses. It is an on-chip bus specification that details a strategy for the interconnection and management of functional blocks that make up a *System-on-Chip* (SoC). It aids in the development of embedded processors with one or more CPUs or signal processors and multiple peripherals. AMBA complements a reusable design methodology by defining a common backbone for SoC modules.

AHB *See* Advanced High-performance Bus.

AHB-AP *See* AHB Access Port.

AMBA *See* Advanced Microcontroller Bus Architecture.

Advanced Trace Bus (ATB)

A bus used by trace devices to pass trace data around a CoreSight system in a trace protocol agnostic manner.

ATB

See Advanced Trace Bus.

AXI

See Advanced eXtensible Interface.

Byte

An 8-bit data item.

CoreSight

The infrastructure for monitoring, tracing, and debugging a complete system on chip.

Data cache

A block of on-chip fast access memory locations, situated between the processor and main memory, used for storing and retrieving copies of often used data. This is done to greatly increase the average speed of memory accesses and so improve processor performance.

Embedded Trace Buffer

The ETB provides on-chip storage of trace data using a configurable sized RAM.

Embedded Trace Macrocell (ETM)

A hardware macrocell that, when connected to a processor core, outputs instruction and data trace information on a trace port. The ETM provides processor driven trace through a trace port compliant to the ATB protocol.

Macrocell

A complex logic block with a defined interface and behavior. A typical VLSI system comprises several macrocells (such as a processor, an ETM, and a memory block) plus application-specific logic.

Microprocessor

See Processor.

Processor

A processor is the circuitry in a computer system required to process data using the computer instructions. It is an abbreviation of microprocessor. A clock source, power supplies, and main memory are also required to create a minimum complete working computer system.

Trace funnel

A device that combines multiple trace sources onto a single bus.

Trace port

A port on a device, such as a processor or ASIC, used to output trace information.

Trace Port Analyzer (TPA)

A hardware device that captures trace information output on a trace port. This can be a low-cost product designed specifically for trace acquisition, or a logic analyzer.

Write

Writes are defined as operations that have the semantics of a store. That is, the ARM instructions SRS, STM, STRD, STC, STRT, STRH, STRB, STRBT, STREX, SWP, and SWPB, and the Thumb instructions STM, STR, STRH, STRB, and PUSH.

Java instructions that are accelerated by hardware can cause a number of writes to occur, according to the state of the Java stack and the implementation of the Java hardware acceleration.